

## Pretty printing—a problem in your daily life

We would like to print a given document nicely, e.g.,

- right ends of lines are justified,
- spaces between words (or letters) are as even as possible,
- the number of words broken into two lines (by using “-”) is as small as possible.

Applications: newspapers, books, word processors, T<sub>E</sub>X, L<sub>A</sub>T<sub>E</sub>X, etc.

## Layout examples:

C'erano una volta tre porcellini fratelli che vivevano con i loro genitori. I tre porcellini mangiavano bevevano e si divertivano in tutti i giorni di sole, mentre in quelli di pioggia si riposavano e con il passare del tempo finirono per crescere ben pasciuti.

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(↖) not right-justified

(↑) right-justified;

line-breaks: the same as (↖)

(←) typeset by L<sup>A</sup>T<sub>E</sub>X

Below we assume for simplicity that it is not allowed to divide a word.

# The algorithm below can easily be generalized to the case where line-breaking in a word is allowed and spacing between characters (even in a word) can vary. In L<sup>A</sup>T<sub>E</sub>X, both of these are considered!

**Input:** The width  $l_i$  of each word  $w_i$  ( $i = 1, 2, \dots, n$ ),

an ideal space width  $s$ , and a text width  $L$ .

When words from  $w_i$  to  $w_j$  ( $1 \leq i \leq j \leq n$ ) are written in a line, the width of each space between words is

$$s(i, j) = \frac{L - \sum_{k=i}^j l_k}{\max\{j - i, 1\}},$$

which incurs a cost

$$c(i, j) = \begin{cases} +\infty, & \text{if } s(i, j) \leq 0 \\ (s - s(i, j))^2, & \text{if } s(i, j) > 0 \text{ and } j < n \\ \max\{s - s(i, j), 0\}^2, & \text{if } s(i, j) > 0 \text{ and } j = n. \end{cases}$$

**Output:** A layout (points to break lines) that minimizes the total cost.

## DP for pretty printing

$f^*(i)$ : the cost of the best layout for words  $w_i, w_{i+1}, \dots, w_n$ .

The recurrence formula for  $f^*(i)$ :

$$f^*(i) = \begin{cases} 0, & i = n + 1 \\ \min_{i \leq j \leq n} \{c(i, j) + f^*(j + 1)\}, & i \leq n. \end{cases}$$

### Computation time of the DP:

- If  $c(i, j)$  is computed from scratch  $\implies O(n^3)$  time in total.
- Compute  $c(i, j)$  in  $O(1)$  time using  $\sum_{k=i}^j l_k = \sum_{k=i}^{j-1} l_k + l_j$   
 $\implies O(n^2)$  time in total.

**Note:** For cost function  $c(i, j)$ , you can use different functions, e.g., cubic power, as long as its computation is easy.